SciencesPo Computational Economics Spring 2019

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1 Optimization 2: Algorithms and Constraints

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1.1 Bracketing

- A derivative-free method for *univariate f*
- works only on **unimodal** *f*
- (Draw choosing initial points and where to move next)

1.2 The Golden Ratio or Bracketing Search for 1D problems

- A derivative-free method
- a Bracketing method
	- **–** find the local minimum of *f* on [*a*, *b*]
	- **–** select 2 interior points *c*, *d* such that $a < c < d < b$
		- *∗* $f(c) ≤ f(d) \implies$ min must lie in [*a*, *d*]. replace *b* with *d*, start again with [*a*, *d*]
		- *∗* $f(c) > f(d)$ \implies min must lie in [*c*, *b*]. replace *a* with *c*, start again with [*c*, *b*]
	- **–** how to choose *b*, *d* though?
	- **–** we want the length of the interval to be independent of whether we replace upper or lower bound
	- **–** we want to reuse the non-replaced point from the previous iteration.
	- **–** these imply the golden rule:
	- **–** new point $x_i = a + \alpha_i(b a)$, where $\alpha_1 = \frac{3 \sqrt{5}}{2}$ $\frac{\sqrt{5}}{2}$, $\alpha_2 =$ *√* 5*−*1 2
	- **–** *α*² is known as the *golden ratio*, well known for it's role in renaissance art.

In [1]: **using** Plots

```
using Optim
gr()
f(x) = exp(x) - x^4minf(x) = -f(x)brent = optimize(min, 0, 2, Brent())golden = optimize(minf,0,2,GoldenSection())
```

```
println("brent = $brent")
       println("golden = $golden")
       plot(f, 0, 2)brent = Results of Optimization Algorithm
 * Algorithm: Brent's Method
* Search Interval: [0.000000, 2.000000]
 * Minimizer: 8.310315e-01
* Minimum: -1.818739e+00
* Iterations: 12
 * Convergence: max(|x - x_upper|, |x - x_lower|) <= 2*(1.5e-08*|x|+2.2e-16): true
 * Objective Function Calls: 13
golden = Results of Optimization Algorithm
* Algorithm: Golden Section Search
 * Search Interval: [0.000000, 2.000000]
 * Minimizer: 8.310315e-01
* Minimum: -1.818739e+00
 * Iterations: 37
```
- * Convergence: max(|x x_upper|, |x x_lower|) <= 2*(1.5e-08*|x|+2.2e-16): true
- * Objective Function Calls: 38

```
Out[1]:
```


1.2.1 Bisection Methods

- Root finding: Roots.jl
- Root finding in multivariate functions: [IntervalRootFinding.jl](https://github.com/JuliaIntervals/IntervalRootFinding.jl/)

```
In [80]: using Roots
         #ăfind the zeros of this function:
        f(x) = exp(x) - x^4## bracketing
        fzero(f, 8, 9) # 8.613169456441398
        fzero(f, -10, 0) # -0.8155534188089606
Out[80]: -0.8155534188089606
In [36]: using IntervalRootFinding, IntervalArithmetic
        -10.10Out[36]: [-10, 10]
In [37]: X = IntervalBox(1..3, 2..4)Out[37]: [1, 3] Œ [2, 4]
In [38]: a = 0interval(0.1, 0.3)b = 0interval(0.3, 0.6)a + bOut[38]: [0.399999, 0.900001]
In [41]: rts = roots(x->x^2 - 2, -10..10, IntervalRootFinding.Bisection)
Out[41]: 2-element Array{Root{Interval{Float64}},1}:
         Root([1.41377, 1.41439], :unknown)
         Root([-1.41471, -1.41407], :unknown)
```
1.3 Rosenbrock Banana and Optim.jl

• We can supply the objective function and - depending on the solution algorithm - the gradient and hessian as well.

```
In [4]: using Optim
        using OptimTestProblems
        for (name, prob) in MultivariateProblems.UnconstrainedProblems.examples
          println(name)
        end
```
Rosenbrock Quadratic Diagonal Hosaki Large Polynomial

Penalty Function I Beale Extended Rosenbrock Polynomial Powell Exponential Paraboloid Diagonal Paraboloid Random Matrix Extended Powell Trigonometric Fletcher-Powell Parabola Himmelblau

In [5]: rosenbrock = MultivariateProblems.UnconstrainedProblems.examples["Rosenbrock"]

Out [5]: OptimTestProblems.MultivariateProblems.OptimizationProblem{Nothing,Nothing,Float64,Str

1.4 Comparison Methods

- We will now look at a first class of algorithms, which are very simple, but sometimes a good starting point.
- They just *compare* function values.
- *Grid Search* : Compute the objective function at $G = \{x_1, \ldots, x_N\}$ and pick the highest value of *f* .
	- **–** This is very slow.
	- **–** It requires large *N*.
	- **–** But it's robust (will find global optimizer for large enough *N*)

```
In [44]: # grid search on rosenbrock
         grid = collect(-1.0:0.1:3);grid2D = [[i;j] for i in grid, j in grid];val2D = map(rosenbrock.f,grid2D);
         r = \text{findmin}(val2D);
         println("grid search results in minimizer = $(grid2D[r[2]])")
```
grid search results in minimizer = [1.0, 1.0]

1.5 Local Descent Methods

- Applicable to multivariate problems
- We are searching for a *local model* that provides some guidance in a certain region of *f* over **where to go to next**.
- Gradient and Hessian are informative about this.

1.5.1 Local Descent Outline

All descent methods follow more or less this structure. At iteration *k*,

- 1. Check if candidate **x** (*k*) satisfies stopping criterion:
	- if yes: stop
	- if no: continue
- 2. Get the local *descent direction* **d** (*k*) , using gradient, hessian, or both.
- 3. Set the $step\ size$, i.e. the length of the next step, α^k
- 4. Get the next candidate via

$$
\mathbf{x}^{(k+1)} \longleftarrow \alpha^k \mathbf{d}^{(k)}
$$

1.5.2 The Line Search Strategy

- An algorithm from the line search class chooses a direction **d** (*k*) *∈* **R***ⁿ* and searches along that direction starting from the current iterate $x_k \in \mathbb{R}^n$ for a new iterate $x_{k+1} \in \mathbb{R}^n$ with a lower function value.
- After deciding on a direction **d** (*k*) , one needs to decide the *step length α* to travel by solving

$$
\min_{\alpha>0} f(x_k + \alpha \mathbf{d}^{(k)})
$$

- In practice, solving this exactly is too costly, so algos usually generate a sequence of trial values *α* and pick the one with the lowest *f* .
- In [46]: *# https://github.com/JuliaNLSolvers/LineSearches.jl*

using LineSearches

```
algo_hz = Optim.Newton(linesearch = HagerZhang()) # Both Optim.jl and IntervalRoot.
res_hz = Optim.optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, rosenbrock.initial
```
- Out[46]: Results of Optimization Algorithm
	- * Algorithm: Newton's Method
	- * Starting Point: [-1.2,1.0]
	- * Minimizer: [1.0000000000000033,1.0000000000000067]
	- * Minimum: 1.109336e-29
	- * Iterations: 23
	- * Convergence: true
		- * |x x'| 0.0e+00: false
		- $|x x'| = 1.13e-08$
		- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 6.35e+13 |f(x)|$
		- * |g(x)| 1.0e-08: true $|g(x)| = 6.66e-15$
		- * Stopped by an increasing objective: false
		- * Reached Maximum Number of Iterations: false
	- * Objective Calls: 71
	- * Gradient Calls: 71
	- * Hessian Calls: 23

1.5.3 The Trust Region Strategy

- First choose max step size, then the direction
- Finds the next step $\bar{\mathbf{x}}^{(k+1)}$ by minimizing a model of \hat{f} over a *trust region*, centered on $\mathbf{x}^{(k)}$

– 2nd order Tayloer approx of *f* is common.

- Radius δ of trust region is changed based on how well \hat{f} fits f in trust region.
- Get **x** *′* via

$$
\min_{\mathbf{x}'} \quad \hat{f}(\mathbf{x}') \n\text{subject to} \quad \|\mathbf{x} - \mathbf{x}' \le \delta\|
$$

In [47]: *# Optim.jl has a TrustRegion for Newton (see below for Newton's Method)*

```
NewtonTrustRegion(; initial_delta = 1.0, # The starting trust region radius
```

```
delta_hat = 100.0, # The largest allowable trust region radius
```

```
eta = 0.1, #When rho is at least eta, accept the step.
```

```
rho_lover = 0.25, # When rho is less than rho_lower, shrink the t
rho_upper = 0.75) # When rho is greater than rho_upper, grow the
```

```
res = Optim.optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, rosenbrock.initial_x
```

```
Out[47]: Results of Optimization Algorithm
```
- * Algorithm: Newton's Method (Trust Region)
- * Starting Point: [-1.2,1.0]
- * Minimizer: [0.9999999994405535,0.9999999988644926]
- * Minimum: 3.405841e-19
- * Iterations: 25
- * Convergence: true
	- $* |x x'|$ 0.0e+00: false
		- $|x x'| = 8.84e-06$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 1.87e+08 |f(x)|$
	- * |g(x)| 1.0e-08: true $|g(x)| = 5.53e-09$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: false
- * Objective Calls: 26
- * Gradient Calls: 26
- * Hessian Calls: 22

1.5.4 Stopping criteria

- 1. maximum number of iterations reached
- 2. absolute improvement $|f(x) f(x')| \leq \epsilon$
- 3. relative improvement $|f(x) f(x')| / |f(x)| \le \epsilon$
- 4. Gradient close to zero $|g(x)| \approx 0$

1.5.5 Gradient Descent

• Here we define

$$
\mathbf{g}^{(k)} = \nabla f(\mathbf{d}^{(k)})
$$

• And our descent becomes

$$
\mathbf{d}^{(k)} = -\nabla \frac{\mathbf{g}^{(k)}}{\|\mathbf{g}^{(k)}\|}
$$

• Minimizing wrt step size results in a jagged path (each direction is orthogonal to previous direction!)

α^(*k*) = arg min *αf*($\mathbf{x}^{(k)} + \alpha \mathbf{d}^{(k)}$)

• *Conjugate* Gradient avoids this issue.

```
In [48]: # Optim.jl again
         GradientDescent(; alphaguess = LineSearches.InitialPrevious(),
                           linesearch = LineSearches.HagerZhang(),
                           P = nothing,precondprep = (P, x) -> nothing)
```

```
Out[48]: GradientDescent{InitialPrevious{Float64},HagerZhang{Float64,Base.RefValue{Bool}},Noth
           alpha: Float64 1.0
           alphamin: Float64 0.0
           alphamax: Float64 Inf
         , HagerZhang{Float64,Base.RefValue{Bool}}
           delta: Float64 0.1
           sigma: Float64 0.9
           alphamax: Float64 Inf
           rho: Float64 5.0
           epsilon: Float64 1.0e-6
           gamma: Float64 0.66
           linesearchmax: Int64 50
           psi3: Float64 0.1
           display: Int64 0
           mayterminate: Base.RefValue{Bool}
         , nothing, getfield(Main, Symbol("##49#50"))(), Flat())
In [49]: # there is a dedicated LineSearch package: https://github.com/JuliaNLSolvers/LineSear
         GD = optimize(rosenbrock.f, rosenbrock.g!,[0.0, 0.0],GradientDescent())
         GD1 = optimize(rosenbrock.f, rosenbrock.g!,[0.0, 0.0],GradientDescent(),Optim.Options
         GD2 = optimize(rosenbrock.f, rosenbrock.g!,[0.0, 0.0],GradientDescent(),Optim.Options
         println("gradient descent = $GD")
         println("\n")
         println("gradient descent 2 = $GD1")
         println("\n")
         println("gradient descent 3 = $GD2")
gradient descent = Results of Optimization Algorithm
 * Algorithm: Gradient Descent
 * Starting Point: [0.0,0.0]
 * Minimizer: [0.9356732500354086,0.875073922357589]
```
- * Minimum: 4.154782e-03
- * Iterations: 1000
- * Convergence: false
	- * |x x'| 0.0e+00: false $|x - x'| = 1.82e-04$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 1.97e-03 |f(x)|$
	- * $|g(x)|$ 1.0e-08: false $|g(x)| = 8.21e-02$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: true
- * Objective Calls: 2532
- * Gradient Calls: 2532

gradient descent 2 = Results of Optimization Algorithm

- * Algorithm: Gradient Descent
- * Starting Point: [0.0,0.0]
- * Minimizer: [0.9978398797724763,0.9956717950747302]
- * Minimum: 4.682073e-06
- * Iterations: 5000
- * Convergence: false
	- * |x x'| 0.0e+00: false $|x - x'| = 5.08e-06$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 1.62e-03 |f(x)|$
	- * $|g(x)|$ 1.0e-08: false $|g(x)| = 2.53e-03$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: true
- * Objective Calls: 12532
- * Gradient Calls: 12532

gradient descent 3 = Results of Optimization Algorithm

- * Algorithm: Gradient Descent
- * Starting Point: [0.0,0.0]
- * Minimizer: [0.9999999914304203,0.9999999828109042]
- * Minimum: 7.368706e-17

```
* Iterations: 20458
```
- * Convergence: true
	- * |x x'| 0.0e+00: false $|x - x'| = 2.00e-11$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 1.61e-03 |f(x)|$
	- $* |g(x)|$ 1.0e-08: true

```
|g(x)| = 9.99e-09
```
- * Stopped by an increasing objective: false
- * Reached Maximum Number of Iterations: false
- * Objective Calls: 51177

* Gradient Calls: 51177

1.6 Second Order Methods

1.6.1 Newton's Method

• We start with a 2nd order Taylor approx over x at step *k*:

$$
q(x) = f(x^{(k)}) + (x - x^{(k)})f'(x^{(k)}) + \frac{(x - x^{(k)})^2}{2}f''(x^{(k)})
$$

• We set find it's root and rearrange to find the next step $k + 1$:

$$
\frac{\partial q(x)}{\partial x} = f'(x^{(k)}) + (x - x^{(k)})f''(x^{(k)}) = 0
$$

$$
x^{(k+1)} = x^{(k)} - \frac{f'(x^{(k)})}{f''(x^{(k)})}
$$

- The same argument works for multidimensional functions by using Hessian and Gradient
- We would get a descent **d**^{*k*} by setting:

$$
\mathbf{d}^k = -\frac{\mathbf{g}^k}{\mathbf{H}^k}
$$

- There are several options to avoid (often costly) computation of the Hessian **H**:
- 1. Quasi-Newton updates **H** starting from identity matrix
- 2. Broyden-Fletcher-Goldfarb-Shanno (BFGS) does better with approx linesearch
- 3. L-BFGS is the limited memory version for large problems

In [6]: optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, [0.0, 0.0], Optim.Newton(),Optim.


```
Out[6]: Results of Optimization Algorithm
```
- * Algorithm: Newton's Method
- * Starting Point: [0.0,0.0]
- * Minimizer: [0.9999999999999994,0.9999999999999989]
- * Minimum: 3.081488e-31
- * Iterations: 14
- * Convergence: true
	- * |x x'| 0.0e+00: false
		- $|x x'| = 3.06e-09$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false $|f(x) - f(x')| = 3.03e+13 |f(x)|$
	- * $|g(x)|$ 1.0e-08: true
	- $|g(x)| = 1.11e-15$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: false
- * Objective Calls: 44
- * Gradient Calls: 44
- * Hessian Calls: 14

In [7]: @show optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, [-1.0, 3.0], BFGS());

optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, [-1.0, 3.0], BFGS()) = Results of Optimiz

```
* Algorithm: BFGS
```
- * Starting Point: [-1.0,3.0]
- * Minimizer: [0.9999999999999956,0.999999999999987]
- * Minimum: 1.707144e-27
- * Iterations: 39
- * Convergence: true
	- * |x x'| 0.0e+00: false
		- $|x x'| = 1.54e-08$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false
	- $|f(x) f(x')| = 3.55e+10 |f(x)|$
	- * |g(x)| 1.0e-08: true $|g(x)| = 1.63e-12$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: false
- * Objective Calls: 137
- * Gradient Calls: 137

In [8]: *# low memory BFGS*

@show optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, [0.0, 0.0], LBFGS());

optimize(rosenbrock.f, rosenbrock.g!, rosenbrock.h!, [0.0, 0.0], LBFGS()) = Results of Optimiz

- * Algorithm: L-BFGS
- * Starting Point: [0.0,0.0]
- * Minimizer: [0.999999999999928,0.9999999999998559]
- * Minimum: 5.191703e-27
- * Iterations: 24
- * Convergence: true
	- * |x x'| 0.0e+00: false $|x - x'| = 4.58e-11$
	- * $|f(x) f(x')|$ 0.0e+00 $|f(x)|$: false
	- $|f(x) f(x')| = 8.50e+07 |f(x)|$
	- * |g(x)| 1.0e-08: true
	- $|g(x)| = 1.44e-13$
	- * Stopped by an increasing objective: false
	- * Reached Maximum Number of Iterations: false
- * Objective Calls: 67
- * Gradient Calls: 67

Direct Methods

- No derivative information is used *derivative free*
- If it's very hard / impossible to provide gradient information, this is our only chance.
- Direct methods use other criteria than the gradient to inform the next step (and ulimtately convergence).

1.6.2 Cyclic Coordinate Descent – Taxicab search

- We do a line search over each dimension, one after the other
- *taxicab* because the path looks like a NYC taxi changing direction at each block.
- given $x^{(1)}$, we proceed

$$
\mathbf{x}^{(2)} = \arg\min_{x_1} f(x_1, x_2^{(1)}, \dots, x_n^{(1)})
$$

$$
\mathbf{x}^{(3)} = \arg\min_{x_2} f(x_1^{(2)}, x_2, x_3^{(2)}, \dots, x_n^{(2)})
$$

• unfortunately this can easily get stuck because it can only move in 2 directions.

```
In [9]: # start to setup a basis function, i.e. unit vectors to index each direction:
        basis(i, n) = [k == i ? 1.0 : 0.0 for k in 1 : n]
        function cyclic_coordinate_descent(f, x, )
            , n = \text{Inf}, length(x)while abs() >
                x = \text{copy}(x)for i in 1 : n
                         d = basis(i, n)x = line\_search(f, x, d)end
                 = norm(x - x)end
            return x
        end
```
Out[9]: cyclic_coordinate_descent (generic function with 1 method)

1.6.3 General Pattern Search

- We search according to an arbitrary *pattern P* of candidate points, anchored at current guess **x**.
- With step size *α* and set *D* of directions

$$
\mathcal{P} = \mathbf{x} + \alpha \mathbf{d} \text{ for } \mathbf{d} \in \mathcal{D}
$$

- Convergence is guaranteed under conditions:
	- **–** *D* must be a positive spanning set: at least one **d** *∈ D* has a non-zero gradient.

```
In [10]: function generalized_pattern_search(f, x, , D, , =0.5)
             y, n = f(x), length(x)
             evals = 0while >
                 improved = false
                 for (i,d) in enumerate(D)
                    x = x + *dy = f(x)evals += 1if y < y
                         x, y, improved = x, y, true
                         D = pushfirst!(deleteat!(D, i), d)break
                     end
                 end
                 if !improved
                      *=end
             end
            println("$evals evaluations")
             return x
```

```
end
```
Out[10]: generalized_pattern_search (generic function with 2 methods)

```
In [11]: D = [[1,0], [0,1], [-1,-0.5]]D = [[1, 0], [0, 1]]y=generalized_pattern_search(rosenbrock.f,zeros(2),0.8,D,1e-6 )
```
11923 evaluations

```
Out[11]: 2-element Array{Float64,1}:
          0.9996734619140493
          0.9993469238280956
```
1.7 Bracketing for Multidimensional Problems: Nelder-Mead

- The Goal here is to find the simplex containing the local minimizer *x ∗*
- In the case where f is n-D, this simplex has $n + 1$ vertices
- In the case where f is 2-D, this simplex has $2 + 1$ vertices, i.e. it's a triangle.
- The method proceeds by evaluating the function at all $n + 1$ vertices, and by replacing the worst function value with a new guess.
- this can be achieved by a sequence of moves:
	- **–** reflect
	- **–** expand
	- **–** contract
	- **–** shrink movements.

- this is a very popular method. The matlab functions fmincon and fminsearch implements it.
- When it works, it works quite fast.
- No derivatives required.

```
In [12]: nm=optimize(rosenbrock.f, [0.0, 0.0], NelderMead());
         nm.minimizer
```

```
Out[12]: 2-element Array{Float64,1}:
          0.9999634355313174
          0.9999315506115275
```
• But.

1.8 Bracketing for Multidimensional Problems: Comment on Nelder-Mead

Lagarias et al. (SIOPT, 1999): At present there is no function in any dimension greater than one, for which the original Nelder-Mead algorithm has been proved to converge to a minimizer.

Given all the known inefficiencies and failures of the Nelder-Mead algorithm [. . .], one might wonder why it is used at all, let alone why it is so extraordinarily popular.

1.9 things to read up on

- Divided Rectangles (DIRECT)
- simulated annealing and other stochastic gradient methods

1.10 Stochastic Optimization Methods

- Gradient based methods like steepest descent may be susceptible to getting stuck at local minima.
- Randomly shocking the value of the descent direction may be a solution to this.
- For example, one could modify our gradient descent from before to become

$$
\mathbf{x}^{(k+1)} \longleftarrow \mathbf{x}^{(k)} + \alpha^k \mathbf{g}^{(k)} + \mathbf{w}^{(k)}
$$

- where $^{\prime\prime}(k) \sim N(0, \sigma_k^2)$, decreasing with *k*.
- This *stochastic gradient descent* is often used when training neural networks.

1.10.1 Simulated Annealing

- We specify a *temperature* that controls the degree of randomness.
- At first the temperature is high, letting the search jump around widely. This is to escape local minima.
- The temperature is gradually decreased, reducing the step sizes. This is to find the local optimimum in the *best* region.
- At every iteration *k*, we accept new point **x** *′* with

$$
Pr(\text{accept } \mathbf{x}') = \begin{cases} 1 & \text{if } \Delta y \le 0 \\ \min(e^{\Delta y/t}, 1) & \text{if } \Delta y > 0 \end{cases}
$$

- here $\Delta y = f(x') f(x)$, and *t* is the *temperature*.
- Pr(accept **x** *′*) is called the **Metropolis Criterion**, building block of *Accept/Reject* algorithms.

```
In [15]: #ăf: function
```

```
# x: initial point
# T: transition distribution
#ăt: temp schedule, k_max: max iterations
function simulated_annealing(f, x, T, t, k_max)
    y = f(x)ytrace = zeros(typeof(y), k_max)
```

```
x_b best, y_b best = x, yfor k in 1 : k_max
        x = x + rand(T)y = f(x)y = y - yif y \cup \lvert \mathbf{rand}() \leq \exp(-y/t(k))x, y = x, yend
         if y < y_best
             x_b best, y_b best = x, yend
        ytrace[k] = y_bestend
    return x_best,ytrace
end
```

```
Out[15]: simulated_annealing (generic function with 1 method)
```

```
In [1]: function ackley(x, a=20, b=0.2, c=2)
            d = length(x)return -a*exp(-b*sqrt(sum(x.^2)/d)) - exp(sum(cos.(c*xi) for xi in x)/d) + a + exp(cos.end
        using Plots
        gr()
        surface(-30:0.1:30,-30:0.1:30,(x,y)->ackley([x, y]),cbar=false)
```

```
Out[1]:
```


```
In [16]: p = Any[]
         using Distributions
         gr()
         nitters = 1000temps = (1, 10, 25)push!(p,[plot(x->i/x,1:1000,title = "tmp $i",lw=2,ylims = (0,1),leg = false) for i in (1,10,25)]...)
         for sig in (1,5,25), t1 in (1,10,25)
             y = simulated_annealing(ackley, [15,15],MvNormal(2,sig),x->t1/x,1000)[2][:]
             push!(p,plot(y,title = "sig = $sig",leg=false,lw=1.5,color="red",ylims = (0,20)))
         end
         plot(p..., layout = (4,3))
```


2 Constraints

Recall our core optimization problem:

$$
\min_{x \in \mathbb{R}^n} f(x) \text{ s.t. } x \in \mathcal{X}
$$

- Up to now, the feasible set was $\mathcal{X} \in \mathbb{R}^n$.
- In **constrained problems** X is a subset thereof.
- We already encountered *box constraints*, e.g. $x \in [a, b]$.
- Sometimes the contrained solution coincides with the unconstrained one, sometimes it does not.
- There are *equality constraints* and *inequality constraints*.

2.1 Lagrange Multipliers

• Used to optimize a function subject to equality constraints.

$$
\min_{x} f(x)
$$

subject to $h(x) = 0$

where both *f* and *h* have continuous partial derivatives.

• We look for contour lines of f that are aligned to contours of $h(x) = 0$.

In other words, we want to find the best *x* s.t. $h(x) = 0$ and we have

$$
\nabla f(x) = \lambda \nabla h(x)
$$

for some *Lagrange Mutliplier λ* * Notice that we need the scalar *λ* because the magnitudes of the gradients may be different. * We therefore form the the **Lagrangian**:

$$
\mathcal{L}(x,\lambda) = f(x) - \lambda h(x)
$$

2.1.1 Example

Suppose we have

$$
\min_{x} - \exp\left(-\left(x_1 x_2 - \frac{3}{2}\right)^2 - \left(x_2 - \frac{3}{2}\right)^2\right)
$$
\nsubject to $x_1 - x_2^2 = 0$

We form the Lagrangiagn:

$$
\mathcal{L}(x_1, x_2, \lambda) = -\exp\left(-\left(x_1x_2 - \frac{3}{2}\right)^2 - \left(x_2 - \frac{3}{2}\right)^2\right) - \lambda(x_1 - x_2^2)
$$

Then we compute the gradient wrt to x_1 , x_2 , λ , set to zero and solve.

In [11]: gr() $f(x1, x2) = -exp.(- (x1. *x2 - 3/2).^2 - (x2-3/2).^2)$ $c(x1) = sqrt(x1)$ x=0:0.01:3.5 contour(x,x,(x,y)->f(x,y),lw=1.5,levels=[collect(0:-0.1:-0.85)...,-0.887,-0.95,-1]) plot!(c,0.01,3.5,label="",lw=2,color=:black) scatter!($[1.358]$, $[1.165]$,markersize=5,markercolor=:red,label="Constr. Optimum")

• If we had multiple constraints (*l*), we'd just add them up to get

$$
\mathcal{L}(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) - \sum_{i=1}^{l} \lambda_i h_i(\mathbf{x})
$$

2.2 Inequality Constraints

Suppose now we had

$$
\underset{\mathbf{x}}{\text{min}}f(\mathbf{x})
$$

subject to
$$
g(\mathbf{x}) \leq 0
$$

which, if the solution lies right on the constraint *boundary*, means that

$$
\nabla f - \mu \nabla g = 0
$$

for some scalar μ - as before.

- In this case, we say the **constraint is active**.
- In the opposite case, i.e. the solution lies **inside** the contrained region, we way the contraint is **inactive**.
- In that case, we are back to an *unconstrained* problem, look for $\nabla f = 0$, and set $\mu = 0$.

```
In [12]: #ăthe blue area shows the FEASIBLE SET
```

```
contour(x,x,(x,y)->f(x,y),lw=1.5,levels=[collect(0:-0.1:-0.85)...,-0.887,-0.95,-1])
plot!(c,0.01,3.5,label="",lw=2,color=:black,fill=(0,0.5,:blue))
scatter!([1.358],[1.165],markersize=5,markercolor=:red,label="Constr. Optimum")
```
Out[12]:


```
In [13]: #ăthe blue area shows the FEASIBLE SET
         #ăNOW THE CONSTRAINT IS INACTIVE OR SLACK!
        c2(x1) = 1+sqrt(x1)contour(x,x,(x,y)->f(x,y),lw=1.5,levels=[collect(0:-0.1:-0.85)...,-0.887,-0.95,-1])
        plot!(c2,0.01,3.5,label="",lw=2,color=:black,fill=(0,0.5,:blue))
        scatter!([1],[1.5],markersize=5,markercolor=:red,label="Unconstr. Optimum")
```
Out[13]:

2.3 Infinity Step

• We could do an **infinite step** to avoid *infeasible points*:

$$
f_{\infty\text{-step}} = \begin{cases} f(\mathbf{x}) & \text{if } g(\mathbf{x}) \le 0\\ \infty & \text{else.} \end{cases}
$$

$$
= f(\mathbf{x}) + \infty (g(\mathbf{x} > 0))
$$

- Unfortunately, this is discontinous and non-differentiable, i.e. hard to handle for algorithms.
- Instead, we use a *linear penalty* $\mu g(\mathbf{x})$ on the objective if the constraint is violated.
- The penalty provides a lower bound to ∞:

$$
\mathcal{L}(\mathbf{x}, \mu) = f(\mathbf{x}) + \mu g(\mathbf{x})
$$

• We can get back the infinite step by maximizing the penalty:

$$
f_{\infty\text{-step}} = \max_{\mu \geq 0} \mathcal{L}(\mathbf{x}, \mu)
$$

• Every infeasible **x** returns ∞ , all others return $f(\mathbf{x})$

2.4 Kuhn-Karush-Tucker (KKT)

• Our problem thus becomes

$$
\min_{\mathbf{x}} \max_{\mu \geq 0} \mathcal{L}(\mathbf{x}, \mu)
$$

- This is called the **primal problem**. Optimizing this requires:
- 1. $g(\mathbf{x}^*) \leq 0$. Point is feasible.
- 2. $\mu \geq 0$. Penalty goes into the right direction. *Dual feasibility*.
- 3. $\mu g(\mathbf{x}^*) = 0$. Feasible point on the boundary has $g(\mathbf{x}) = 0$, otherwise $g(\mathbf{x}) < 0$ and $\mu = 0$.
- 4. $\nabla f(\mathbf{x}^*) \mu \nabla g(\mathbf{x}^*) = 0$. With an active constraint, we want parallel contours of objective and constraint. When inactive, our optimum just has $\nabla f(\mathbf{x}^*) = 0$, which means $\mu = 0$.

The preceding four conditions are called the Kuhn-Karush-Tucker (KKT) conditions. In the above order, and in general terms, they are:

- 1. Feasibility
- 2. Dual Feasibility
- 3. Complementary Slackness
- 4. Stationarity.

The KKT conditions are the FONCs for problems with smooth constraints.

2.5 Duality

We can combine equality and inequality constraints:

$$
\mathcal{L}(\mathbf{x}, \tilde{\boldsymbol{\cdot}}, \tilde{\boldsymbol{\cdot}}) = f(\mathbf{x}) + \sum_{i} \lambda_i h_i(\mathbf{x}) + \sum_{j} \mu_j g_j(\mathbf{x})
$$

where, notice, we reverted the sign of λ since this is unrestricted.

• The Primal problem is identical to the original problem and just as difficult to solve:

$$
\min_{\mathbf{x}} \max_{\mathbf{y}} \mathcal{L}(\mathbf{x}, \mathbf{y})
$$

• The Dual problem reverses min and max:

$$
\max_{\substack{\sim\\ geq 0, \quad x}} \min_{x} \mathcal{L}(x, \bar{f}, \bar{f})
$$

2.5.1 Dual Values

• The *max-min-inequality* states that for any function $f(a, b)$

$$
\max_{\mathbf{a}}\min_{\mathbf{b}}f(\mathbf{a},\mathbf{b})\leq \min_{\mathbf{b}}\max_{\mathbf{a}}f(\mathbf{a},\mathbf{b})
$$

- Hence, the solution to the dual is a lower bound to the solution of the primal problem.
- The solution to the *dual function*, $\min_{x} \mathcal{L}(x, \bar{z}, \bar{z})$ is the min of a collection of linear functions, and thus always concave.
- It is easy to optimize this.
- In general, solving the dual is easy whenever minimizing *L* wrt *x* is easy.

Penalty Methods

- We can convert the constrained problem back to unconstrained by adding penalty terms for constraint violoations.
- A simple method could just count the number of violations:

$$
p_{\text{count}}(\mathbf{x}) = \sum_i (h_i(\mathbf{x}) \neq 0) + \sum_j (g_j(\mathbf{x} > 0))
$$

• and add this to the objective in an *unconstrained* problem with penalty *ρ >* 0

$$
\min_{\mathbf{x}} f(\mathbf{x}) + \rho p_{\text{count}}(\mathbf{x})
$$

- One can choose the penalty function: for example, a quadratic penaly will produce a smooth objective function
- Notice that *ρ* needs to become very large sometimes here.

Augmented Lagrange Method

• This is very similar, but specific to equality constraints.

Interior Point Method

- Also called *barrier method*.
- These methods make sure that the search point remains always feasible.
- As one approaches the constraint boundary, the barrier function goes to infinity. Properties:
- 1. $p_{\text{barrier}}(\mathbf{x})$ is continuous
- 2. $p_{barrier}(x)$ is non negative
- 3. $p_{barrier}(x)$ goes to infinitey as one approaches the constraint boundary

2.5.2 Barriers

• Inverse Barrier

$$
p_{barrier}(\mathbf{x}) = -\sum_{i} \frac{1}{g_i(\mathbf{x})}
$$

• Log Barrier

$$
p_{\text{barrier}}(\mathbf{x}) = -\sum_{i} \begin{cases} \log(-g_i(\mathbf{x})) & \text{if } g_i(\mathbf{x}) \ge -1\\ 0 & \text{else.} \end{cases}
$$

• The approach is as before, one transforms the problem to an unconstrained one and increases *ρ* until convergence:

$$
\min_{\mathbf{x}} f(\mathbf{x}) + \frac{1}{\rho} p_{barrier}(\mathbf{x})
$$

2.5.3 Examples

$$
\min_{x \in \mathbb{R}^2} \sqrt{x_2} \text{ subject to } \sum_{\substack{x_2 \ge (a_1 x_1 + b_1)^3 \\ x_2 \ge (a_2 x_1 + b_2)^3}}
$$

2.6 Constrained Optimisation with [NLopt.jl](https://github.com/JuliaOpt/NLopt.jl)

- We need to specify one function for each objective and constraint.
- Both of those functions need to compute the function value (i.e. objective or constraint) *and* it's respective gradient.
- NLopt expects contraints **always** to be formulated in the format

$$
g(x)\leq 0
$$

where *g* is your constraint function

- The constraint function is formulated for each constraint at *x*. it returns a number (the value of the constraint at *x*), and it fills out the gradient vector, which is the partial derivative of the current constraint wrt *x*.
- There is also the option to have vector valued constraints, see the documentation.
- We set this up as follows:

```
In [9]: using NLopt
```

```
count = 0 # keep track of # function evaluations
function myfunc(x::Vector, grad::Vector)
    if length(grad) > 0
       grad[1] = 0
```

```
grad[2] = 0.5/sqrt(x[2])end
            global count
            count::Int += 1
            println("f_$count($x)")
            sqrt(x[2])end
        function myconstraint(x::Vector, grad::Vector, a, b)
            if length(grad) > 0
                grad[1] = 3a * (a*x[1] + b)^2grad[2] = -1end
            (a*x[1] + b)^3 - x[2]end
        opt = Opt(:LD_MMA, 2)lower_bounds!(opt, [-Inf, 0.])
        xtol_rel!(opt,1e-4)
        min_objective!(opt, myfunc)
        inequality_constraint!(\text{opt}, (x,g) \rightarrow myconstraint(x,g,2,0), 1e-8)inequality_constraint!(opt, (x,g) \rightarrow myconstant(x,g,-1,1), 1e-8)
        (minfunc, minus,ret) = NLopt.optimize(opt, [1.234, 5.678])println("got $minfunc at $minx after $count iterations (returned $ret)")
f_1([1.234, 5.678])
f_2([0.878739, 5.55137])
f_3([0.826216, 5.0439])
f_4([0.473944, 4.07677])
f_5([0.353898, 3.03085])
f_6([0.333873, 1.97179])
f_7([0.333334, 1.04509])
f_8([0.333334, 0.469503])
f_9([0.333333, 0.305792])
f_10([0.333333, 0.296322])
f_11([0.333333, 0.296296])
got 0.5443310477213124 at [0.333333, 0.296296] after 11 iterations (returned XTOL_REACHED)
```
WARNING: using NLopt.optimize! in module Main conflicts with an existing identifier.

2.7 NLopt: Rosenbrock

• Let's tackle the rosenbrock example again.

• To make it more interesting, let's add an inequality constraint.

$$
\min_{x \in \mathbb{R}^2} (1 - x_1)^2 + 100(x_2 - x_1^2)^2
$$
 subject to $0.8 - x_1^2 - x_2^2 \ge 0$

• in NLopt format, the constraint is $x_1 + x_2 - 0.8 \leq 0$

```
In [9]: function rosenbrockf(x::Vector,grad::Vector)
             if length(grad) > 0
                     grad[1] = -2.0 * (1.0 - x[1]) - 400.0 * (x[2] - x[1]^2) * x[1]
                     grad[2] = 200.0 * (x[2] - x[1]^{2})end
             return (1.0 - x[1])^2 + 100.0 * (x[2] - x[1]^2)^2end
        function r_constraint(x::Vector, grad::Vector)
             if length(\text{grad}) > 0
                 grad[1] = 2*x[1]grad[2] = 2*x[2]end
                 return x[1]<sup>2</sup> + x[2]<sup>2</sup> - 0.8
        end
        opt = Opt(:LD_MMA, 2)lower_bounds!(opt, [-5, -5.0])
        min\_objective! (opt, (x,g) \rightarrow rosenbrockf(x,g))inequality_constraint!(opt, (x,g) \rightarrow r_{constant}(x,g))
        ftol_rel!(opt,1e-9)
        NLopt.optimize(opt, [-1.0,0.0])
Out[9]: (0.07588358473630112, [0.724702, 0.524221], :FTOL_REACHED)
```
2.8 JuMP.jl

- Introduce JuMP. j1
- JuMP is a mathematical programming interface for Julia. It is like AMPL, but for free and with a decent programming language.
- The main highlights are:
	- **–** It uses automatic differentiation to compute derivatives from your expression.
	- **–** It supplies this information, as well as the sparsity structure of the Hessian to your preferred solver.
	- **–** It decouples your problem completely from the type of solver you are using. This is great, since you don't have to worry about different solvers having different interfaces.
	- **–** In order to achieve this, JuMP uses [MathProgBase.jl](https://github.com/JuliaOpt/MathProgBase.jl), which converts your problem formulation into a standard representation of an optimization problem.
- Let's look at the readme
- The technical citation is Lubin et al [**?**]

2.9 JuMP: Quick start guide

- this is form the [quick start guide](http://www.juliaopt.org/JuMP.jl/v0.19.0/)
- please check the docs, they are excellent.

2.9.1 Step 1: create a model

- A model collects variables, objective function and constraints.
- it defines a specific solver to be used.
- JuMP makes it very easy to [swap out solver backends](http://www.juliaopt.org/JuMP.jl/dev/installation/) This is very valuable!

```
In [18]: using JuMP
```

```
using GLPK
model = Model(with_optimizer(GLPK.Optimizer))
@variable(model, 0 \leq x \leq 2)\texttt{Quariable}(\texttt{model}, 0 \leq y \leq 30)# next, we set an objective function
Cobjective(model, Max, 5x + 3 * y)
# maybe add a constraint called "con":
@constraint(model, con, 1x + 5y \leq 3);
```
- At this stage JuMP has a mathematical representation of our model internalized
- The MathProgBase machinery knows now exactly how to translate that to different solver interfaces
- For us the only thing left: hit the button!

```
In [15]: JuMP.optimize!(model)
```

```
# look at status
         termination_status(model)
Out[15]: OPTIMAL::TerminationStatusCode = 1
In [16]: # we query objective value and solutions
         @show objective_value(model)
         @show value(x)
         @show value(y)
         # as well as the value of the dual variable on the constraint
         @show dual(con);
objective_value(model) = 10.6
value(x) = 2.0value(y) = 0.2dual(con) = -0.6
```
- The last call gets the *dual value associated with a constraint*
- Economists most of the time call that the *value of the lagrange multiplier*.

For linear programs, a feasible dual on $a \geq$ constraint is nonnegative and a feasible dual on a <= constraint is nonpositive

• This is different to some textbooks and has nothing to do with wether max or minimizing.

In [71]: *# helpfully, we have this, which is always positive:* shadow_price(con)

Out[71]: 0.6

2.10 JuMP handles. . .

- linear programming
- mixed-integer programming
- second-order conic programming
- semidefinite programming, and
- nonlinear programming

```
In [17]: # JuMP: nonlinear Rosenbrock Example
```

```
# Instead of hand-coding first and second derivatives, you only have to give `JuMP` e
# Here is an example.
```

```
using Ipopt
```
let

```
m = Model(with_optimizer(Ipopt.Optimizer))
@variable(m, x)
@variable(m, y)
@NLobjective(m, Min, (1-x)^2 + 100(y-x^2)^2)JuMP.optimize!(m)
@show value(x)
@show value(y)
@show termination_status(m)
```
end

This is Ipopt version 3.12.10, running with linear solver mumps. NOTE: Other linear solvers might be more efficient (see Ipopt documentation).

inequality constraints with lower and upper bounds: 0 inequality constraints with only upper bounds: 0

Number of Iterations...: 14

(scaled) (unscaled) Objective...: 1.3288608467480825e-28 1.3288608467480825e-28 Dual infeasibility...: 2.0183854587685121e-13 2.0183854587685121e-13 Constraint violation...: 0.0000000000000000e+00 0.0000000000000000e+00 Complementarity...: 0.0000000000000000e+00 0.0000000000000000e+00 Overall NLP error...: 2.0183854587685121e-13 2.0183854587685121e-13

```
Number of objective function evaluations = 36
Number of objective gradient evaluations = 15
Number of equality constraint evaluations = 0Number of inequality constraint evaluations = 0Number of equality constraint Jacobian evaluations = 0Number of inequality constraint Jacobian evaluations = 0
Number of Lagrangian Hessian evaluations = 14
Total CPU secs in IPOPT (w/o function evaluations) = 0.006
Total CPU secs in NLP function evaluations = 0.000
```
EXIT: Optimal Solution Found. value(x) = 0.9999999999999899 value(y) = 0.9999999999999792 termination_status(m) = LOCALLY_SOLVED::TerminationStatusCode = 4

Out[17]: LOCALLY_SOLVED::TerminationStatusCode = 4

```
In [18]: # not bad, right?
         # adding the constraint from before:
         let
             m = Model(with_optimizer(Ipopt.Optimizer))
             @variable(m, x)
             @variable(m, y)
             @NLobjective(m, Min, (1-x)^2 + 100(y-x^2)^2)\mathbb{Q}NLconstraint(m, x^2 + y^2 \leq 0.8)JuMP.optimize!(m)
             @show value(x)
             @show value(y)
             @show termination_status(m)
```

```
end
```
This is Ipopt version 3.12.10, running with linear solver mumps. NOTE: Other linear solvers might be more efficient (see Ipopt documentation).

8 7.8681242e-02 0.00e+00 2.12e-03 -2.5 1.12e-02 - 1.00e+00 1.00e+00h 1 9 7.6095770e-02 0.00e+00 6.16e-03 -3.8 1.37e-02 - 1.00e+00 1.00e+00h 1 iter objective inf_pr inf_du lg(mu) ||d|| lg(rg) alpha_du alpha_pr ls 10 7.6033892e-02 0.00e+00 2.23e-06 -3.8 3.99e-04 - 1.00e+00 1.00e+00h 1 11 7.5885642e-02 0.00e+00 2.07e-05 -5.7 7.99e-04 - 1.00e+00 1.00e+00h 1 12 7.5885428e-02 0.00e+00 2.74e-11 -5.7 1.38e-06 - 1.00e+00 1.00e+00h 1 13 7.5883585e-02 0.00e+00 3.19e-09 -8.6 9.93e-06 - 1.00e+00 1.00e+00f 1

Number of Iterations...: 13

(scaled) (unscaled) Objective...: 7.5883585442440671e-02 7.5883585442440671e-02 Dual infeasibility...: 3.1949178858070582e-09 3.1949178858070582e-09 Constraint violation...: 0.0000000000000000e+00 0.0000000000000000e+00 Complementarity...: 2.5454985882932001e-09 2.5454985882932001e-09 Overall NLP error...: 3.1949178858070582e-09 3.1949178858070582e-09

EXIT: Optimal Solution Found. value(x) = 0.7247018392092258 $value(y) = 0.5242206029480763$ termination_status(m) = LOCALLY_SOLVED::TerminationStatusCode = 4

Out[18]: LOCALLY_SOLVED::TerminationStatusCode = 4

2.11 JuMP: Maximium Likelihood

- Let's redo the maximum likelihood example in JuMP.
- Let μ , σ^2 be the unknown mean and variance of a random sample generated from the normal distribution.
- Find the maximum likelihood estimator for those parameters!
- density:

$$
f(x_i|\mu, \sigma^2) = \frac{1}{\sigma\sqrt{2\pi}} \exp\left(-\frac{(x_i - \mu)^2}{2\sigma^2}\right)
$$

• Likelihood Function

$$
L(\mu, \sigma^2) = \Pi_{i=1}^N f(x_i | \mu, \sigma^2) = \frac{1}{(\sigma \sqrt{2\pi})^n} \exp\left(-\frac{1}{2\sigma^2} \sum_{i=1}^N (x_i - \mu)^2\right)
$$

= $(\sigma^2 2\pi)^{-\frac{n}{2}} \exp\left(-\frac{1}{2\sigma^2} \sum_{i=1}^N (x_i - \mu)^2\right)$

- Constraints: *µ ∈* **R**, *σ >* 0
- log-likelihood:

$$
\log L = l = -\frac{n}{2} \log (2\pi\sigma^2) - \frac{1}{2\sigma^2} \sum_{i=1}^{N} (x_i - \mu)^2
$$

• Let's do this in JuMP.

```
In [5]: # Copyright 2015, Iain Dunning, Joey Huchette, Miles Lubin, and contributors
        # example modified
       using Distributions
```

```
let
```

```
distrib = Normal(4.5, 3.5)n = 10000data = rand(distrib,n);m = Model(with_optimizer(Ipopt.Optimizer))
            Quariable(m, mu, start = 0.0)\texttt{Quariable}(m, \text{sigma} \geq 0.0, \text{start} = 1.0)@NLobjective(m, Max, -(n/2)*log(2*sigma^2)-sum((data[i] - mu) ^ 2 for i = 1:n)/(2*s)JuMP.optimize!(m)
            @show termination_status(m)
            printhn(" = ", value(mu), ", mean(data) = ", mean(data))println("2 =", value(sigma)2, ", var(data) = ", var(data))
        end
This is Ipopt version 3.12.10, running with linear solver mumps.
NOTE: Other linear solvers might be more efficient (see Ipopt documentation).
```

```
Number of nonzeros in equality constraint Jacobian...: 0
Number of nonzeros in inequality constraint Jacobian.: 0
Number of nonzeros in Lagrangian Hessian...: 3
Total number of variables...: 2
                  variables with only lower bounds: 1
             variables with lower and upper bounds: 0
```


Number of Iterations...: 60

(scaled) (unscaled) Objective...: 8.5074714157088174e+00 2.6690648196634054e+04 Dual infeasibility...: 9.7569600551044561e-09 3.0610692128632544e-05 Constraint violation...: 0.0000000000000000e+00 0.0000000000000000e+00 Complementarity...: 9.0909090949984745e-10 2.8521078071934833e-06 Overall NLP error...: 9.7569600551044561e-09 3.0610692128632544e-05


```
EXIT: Optimal Solution Found.
termination_status(m) = LOCALLY_SOLVED::TerminationStatusCode = 4
 =
WARNING: using Distributions.mode in module Main conflicts with an existing identifier.
```

```
4.493062898788303, mean(data) = 4.493062936089179
ˆ2 = 12.185571328763903, var(data) = 12.186789996356879
```
3 Linear Constrained Problems (LPs)

• Very similar to before, just that both objective and constraints are *linear*.

$$
\begin{aligned}\n\min_{\mathbf{x}} & \mathbf{r} \mathbf{x} \\
\text{subject to } & \mathbf{w}_{LE}^{(i)T} \mathbf{x} \le b_i \text{ for } i \in 1, 2, 3, \dots \\
&\mathbf{w}_{GE}^{(j)T} \mathbf{x} \ge b_j \text{ for } j \in 1, 2, 3, \dots \\
&\mathbf{w}_{EQ}^{(k)T} \mathbf{x} = b_k \text{ for } k \in 1, 2, 3, \dots\n\end{aligned}
$$

• Our initial JuMP example was of that sort.

3.0.1 Standard Form

- Usually LPs are given in *standard form*
- All constraints are less-than inequalities
- All choice variables are non-negative.

 $\min_{\mathbf{x}} \mathbf{c}^T \mathbf{x}$

subject to $Ax \leq b$

$$
\mathbf{x}\geq 0
$$

- Greater-than inequality constraints are inverted
- equality constraints are split into two
- $x = x^+ x^-$ and we constrain both components to be positive.

3.0.2 Equality Form

$$
\begin{aligned}\n&\text{min}\mathbf{c}^T \mathbf{x} \\
&\text{subject to } \mathbf{A}\mathbf{x} = b \\
&\mathbf{x} \ge 0\n\end{aligned}
$$

• Can transform standard into equality form

$$
\mathbf{A}\mathbf{x} \leq b \to \mathbf{A}\mathbf{x} + \mathbf{s} = b, \mathbf{s} \geq 0
$$

- equality constraints are split into two
- $x = x^+ x^-$ and we constrain both components to be positive.

3.0.3 Solving LPs

- Simplex Algorithm operates on Equality Form
- Moving from one vertex to the next of the feasible set, this is guaranteed to find the optimal solution if the problem is bounded.

3.1 A Cannery Problem

- A can factory (a cannery) has plants in Seattle and San Diego
- They need to decide how to serve markets New-York, Chicago, Topeka
- Firm wants to
	- 1. minimize shipping costs
	- 2. shipments cannot exceed capacity
	- 3. shipments must satisfy demand
- Formalize that!
- Plant capacity cap_i , demands d_j and transport costs from plant *i* to market *j*, $dist_{i,j}c$ are all given.
- Let **x** be a matrix with element *xi*,*^j* for number of cans shipped from *i* to *j*.

3.2 From Maths . . .

$$
\min_{\mathbf{x}} \sum_{i=1}^{2} \sum_{j=1}^{3} dist_{i,j} c \times x_{i,j}
$$
\nsubject to
$$
\sum_{j=1}^{3} x(i,j) \leq cap_i, \forall i
$$
\n
$$
\sum_{i=1}^{2} x(i,j) \geq d_j, \forall j
$$

In [7]: *# ... to JuMP*

- *# https://github.com/JuliaOpt/JuMP.jl/blob/release-0.19/examples/cannery.jl*
- *# Copyright 2017, Iain Dunning, Joey Huchette, Miles Lubin, and contributors*
- *# This Source Code Form is subject to the terms of the Mozilla Public*
- *# License, v. 2.0. If a copy of the MPL was not distributed with this*
- *# file, You can obtain one at http://mozilla.org/MPL/2.0/.*

JuMP

- *# An algebraic modeling language for Julia*
- *# See http://github.com/JuliaOpt/JuMP.jl*

###

```
using JuMP, GLPK, Test
const MOI = JuMP.MathOptInterface
"" ""
    example cannery(; verbose = true)JuMP implementation of the cannery problem from Dantzig, Linear Programming and
Extensions, Princeton University Press, Princeton, NJ, 1963.
Author: Louis Luangkesorn
Date: January 30, 2015
""""
function example_cannery(; verbose = true)
    plants = ["Seattle", "San-Diego"]
    markets = ["New-York", "Chicago", "Topeka"]
    # Capacity and demand in cases.
    capacity = [350, 600]demand = [300, 300, 300]
    # Distance in thousand miles.
    distance = [2.5 1.7 1.8; 2.5 1.8 1.4]
    # Cost per case per thousand miles.
    freight = 90num_plants = length(plants)
    num_markets = length(markets)
    cannery = Model(with_optimizer(GLPK.Optimizer))
    \texttt{Quariable}( \texttt{cannery}, \texttt{ship[1:num_plants, 1:num_markets}] \geq 0)# Ship no more than plant capacity
    @constraint(cannery, capacity_con[i in 1:num_plants],
        sum(ship[i,j] for j in 1:num_markets) <= capacity[i]
    )
    # Ship at least market demand
    @constraint(cannery, demand_con[j in 1:num_markets],
        sum(ship[i,j] for i in 1:num_plants) >= demand[j]
    )
    # Minimize transporatation cost
    @objective(cannery, Min, sum(distance[i, j] * freight * ship[i, j]
        for i in 1:num_plants, j in 1:num_markets)
    \lambdaJuMP.optimize!(cannery)
```

```
if verbose
                println("RESULTS:")
                for i in 1:num_plants
                    for j in 1:num_markets
                        println(" $(plants[i]) $(markets[j]) = $(JuMP.value(ship[i, j]))")
                    end
                end
            end
            @assert JuMP.termination_status(cannery) == MOI.OPTIMAL
            @assert JuMP.primal_status(cannery) == MOI.FEASIBLE_POINT
            @assert JuMP.objective_value(cannery) == 151200.0
        end
        example_cannery()
RESULTS:
  Seattle New-York = 50.0
  Seattle Chicago = 300.0
 Seattle Topeka = 0.0
  San-Diego New-York = 250.0
  San-Diego Chicago = 0.0
  San-Diego Topeka = 300.0
```
4 Discrete Optimization / Integer Programming

- Here the choice variable is contrained to come from a discrete set \mathcal{X} .
- If this set is $X = N$, we have an **integer program**
- If only *some x* have to be discrete, this is a **mixed integer program**

4.1 Example

$$
\min_{\mathbf{x}} x_1 + x_2
$$

subject to $||\mathbf{x}|| \le 2$
 $\mathbf{x} \in \mathbb{N}$

- continuous optimum is (*− √* 2, *− √* 2) and objective is *y* = *−*2 *√* 2 = *−*2.828
- Integer constrained problem is only delivering $y = -2$, and $x^* \in (-2,0), (-1, -1), (0, -2)$

```
In [8]: x = -3:0.01:3
```

```
dx = repeat(range(-3,stop = 3, length = 7),1,7)
contourf(x,x,(x,y)->x+y,color=:blues)scatter!(dx,dx',legend=false,markercolor=:white)
plot!(x->sqrt(4-x^2),-2,2,c=:white)plot!(x->sqrt(4-x^2),-2,2,c=:white)scatter!([-2,-1,0], [0,-1,-2], c=:red)
scatter!([-sqrt(2)],[-sqrt(2)],c=:red,markershape=:cross,markersize=9)
```


4.2 Rounding

- One solution is to just *round the continuous solution to the nearest integer*
- We compute the **relaxed** problem, i.e. the one where *x* is continuous.
- Then we round up or down.
- Can go terribly wrong.

4.3 Cutting Planes

- This is an exact method
- We solve the relaxed problem first.
- Then we add linear constraints that result in the solution becoming integral.

4.4 Branch and Bound

- This enumerates all possible soultions.
- Branch and bound does this, without having to compute all of them.

4.5 Example: The Knapsack Problem

- We are packing our knapsack for a trip but only have space for the most valuable items.
- We have $x_i = 0$ if item *i* is not in the sack, 1 else.

$$
\min_{x} - \sum_{i=1}^{n} v_i x_i
$$

s.t.
$$
\sum_{i=1}^{n} w_i x_i \leq w_{max}
$$

$$
w_i \in \mathbb{N}_{+,v_i} \in \mathbb{R}
$$

- If ther are *n* items, we have 2^n possible design vectors.
- But there is a useful recursive relationship.
- If we solved *n −* 1 knapsack problems so far and deliberate about item *n*
	- **–** If it's not worth including item *n*, then the solution **is** the knapsack problem for *n −* 1 items and capacity *w*max
	- **–** If it IS worth including it: solution will have value of knapsack with *n −* 1 items and reduced capacity, plus the value of the new item
- This **is** dynamic progamming.

4.5.1 Knacksack Recursion

• In particular, the recursion looks like this:

$$
\begin{aligned} &\text{knapsack}\left(i,w_{\max}\right)=\begin{cases} 0 & \text{if } i=0\\ \text{knapsack}\left(i-1,w_{\max}\right) & \text{(discard new item)}\\ \text{max}\begin{cases} \text{knapsack}\left(i-1,w_{\max}-w_i\right)+v_i & \text{(include new item)}\\ \text{knapsack}\left(i-1,w_{\max}-w_i\right)+v_i & \text{(include new item)} \end{cases} & \text{else.}\\ &\text{In } \text{[6]: } \# \text{ Copyright 2017, } \text{Iain Dunning, } \text{Joey Huchette, Miles Lubin, and contributions}\\ &\# \text{ This Source Code Form is subject to the terms of the Mozilla Public} \\ &\text{Hices, v. 2.0. If a copy of the MPL was not distributed with this}\\ &\# \text{ file, You can obtain one at http://mczilla.org/MPL/2.0/.}\\ &\text{http://github.com/JuliaOpt/JwMP.jl}\\ &\# \text{ An algebraic modeling language for Julia}\\ &\# \text{See http://github.com/JuliaOpt/JwMP.jl}\\ &\# \text{knapsack.jl}\\ &\# \text{snams} & \text{sum}(p_j x_j) <= C\\ &\# \text{ so in } \text{array}\\ &\# \text{sum}(w_j x_j) <= C\\ &\# \text{ so in } \text{array}\\ &\# \text{if } \text{sum}(w_j x_j) <= C\\ &\# \text{bin} \# \text{matrix} \text{if } \text{min} \# \text
$$

using JuMP, Cbc, LinearAlgebra

```
# Maximization problem
           m = Model(with_optimizer(Cbc.Optimizer))
            @variable(m, x[1:5], Bin)
            profit = [ 5, 3, 2, 7, 4 ]
            weight = [2, 8, 4, 2, 5]capacity = 10# Objective: maximize profit
            @objective(m, Max, dot(profit, x))
            # Constraint: can carry all
            @constraint(m, dot(weight, x) \leq capacity)# Solve problem using MIP solver
            optimize!(m)
           println("Objective is: ", JuMP.objective_value(m))
           println("Solution is:")
            for i = 1:5print("x[\$i] = ", JuMP.value(x[i]))println(", p[$i]/w[$i] = ", profit[i]/weight[i])
            end
        end
Objective is: 16.0
Solution is:
x[1] = 1.0, p[1]/w[1] = 2.5x[2] = 0.0, p[2]/w[2] = 0.375x[3] = 0.0, p[3]/w[3] = 0.5x[4] = 1.0, p[4]/w[4] = 3.5x[5] = 1.0, p[5]/w[5] = 0.8Welcome to the CBC MILP Solver
Version: 2.9.9
Build Date: Dec 31 2018
command line - Cbc_C_Interface -solve -quit (default strategy 1)
Continuous objective value is 16.5 - 0.00 seconds
Cgl0004I processed model has 1 rows, 5 columns (5 integer (5 of which binary)) and 5 elements
Cutoff increment increased from 1e-05 to 0.9999
Cbc0038I Initial state - 1 integers unsatisfied sum - 0.25
Cbc0038I Solution found of -16
Cbc0038I Before mini branch and bound, 4 integers at bound fixed and 0 continuous
Cbc0038I Mini branch and bound did not improve solution (0.00 seconds)
Cbc0038I After 0.00 seconds - Feasibility pump exiting with objective of -16 - took 0.00 seconds
```

```
let
```
Cbc0012I Integer solution of -16 found by feasibility pump after 0 iterations and 0 nodes (0.0 Cbc0001I Search completed - best objective -16, took 1 iterations and 0 nodes (0.00 seconds) Cbc0035I Maximum depth 0, 4 variables fixed on reduced cost Cuts at root node changed objective from -16.5 to -16 Probing was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cut Gomory was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts Knapsack was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cu Clique was tried 0 times and created 0 cuts of which 0 were active after adding rounds of cuts MixedIntegerRounding2 was tried 0 times and created 0 cuts of which 0 were active after adding FlowCover was tried 0 times and created 0 cuts of which 0 were active after adding rounds of c TwoMirCuts was tried 0 times and created 0 cuts of which 0 were active after adding rounds of

Result - Optimal solution found

In []: